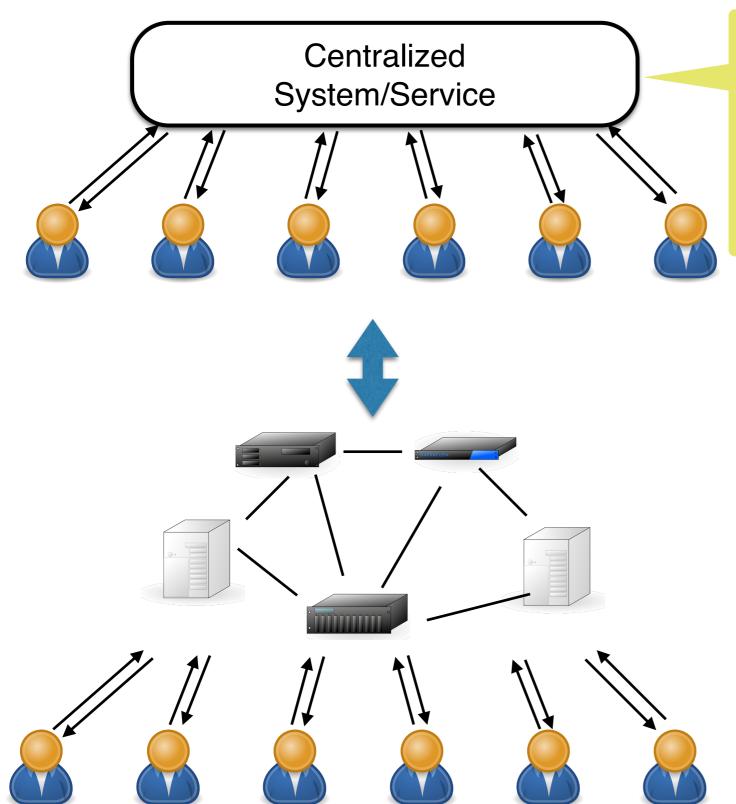


# Secure Decentralization for a Global-Scale Trustworthy Infrastructure

Vassilis Zikas

9.1.2023 NTUA

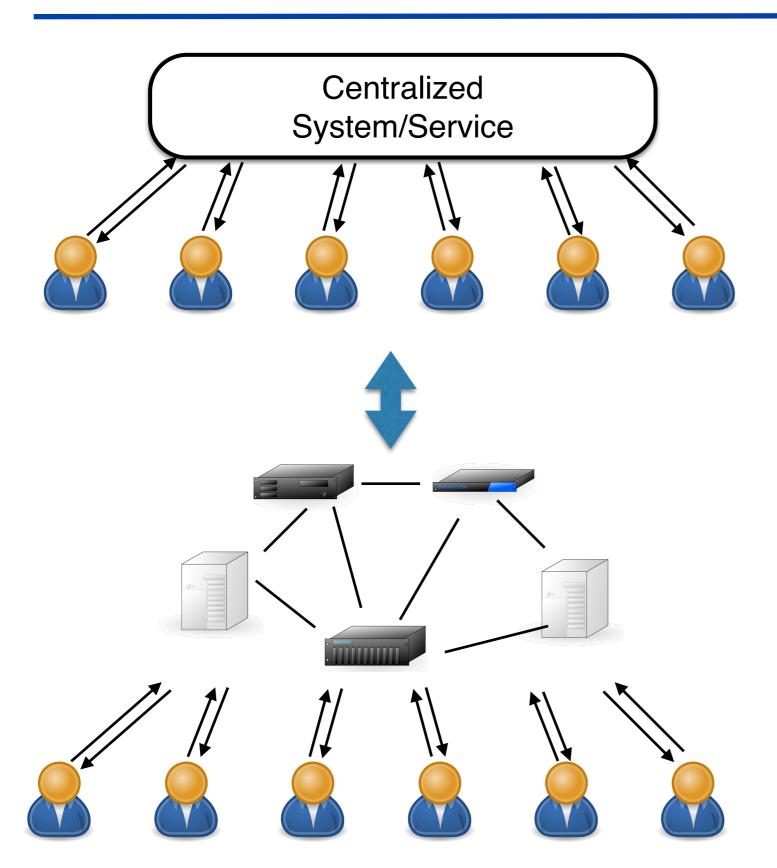
## **The Decentralization Paradigm**



#### Classical Examples:

- Database
- Voting System
- Auctions System
- Authentication Service

## **The Decentralization Paradigm**



#### Why?

- Trust and Power
- Reliability
- Single point of failure (privacy, DoS ..)
- Load (scalability)

## Blockchain as a Scientist/Engineer

#### **Blockchains and Decentralized Ledger Technology (DLT)**

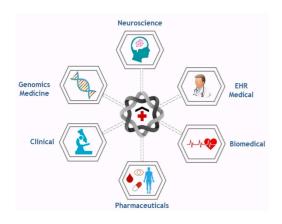






#### Game-changer: Promises to disrupt/improve basic infrastructure:

- Financial Transactions (Cryptocurrencies)
- Trustworthy online computing (Smart contracts)
- Supply chain coordination and tracking
- Fair exchange of digital goods
- Coordinated IoTs
- Biomedical data sharing and access

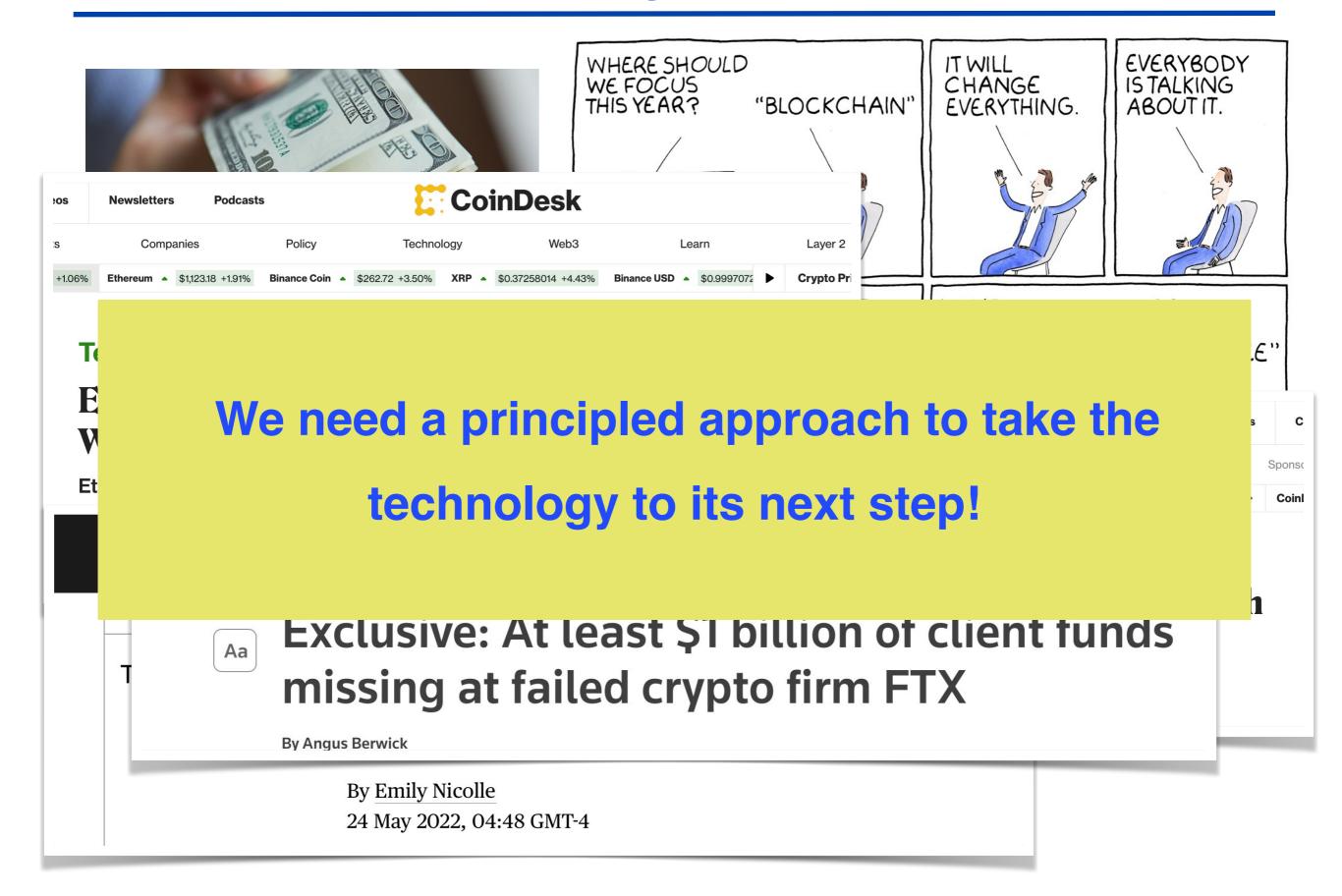








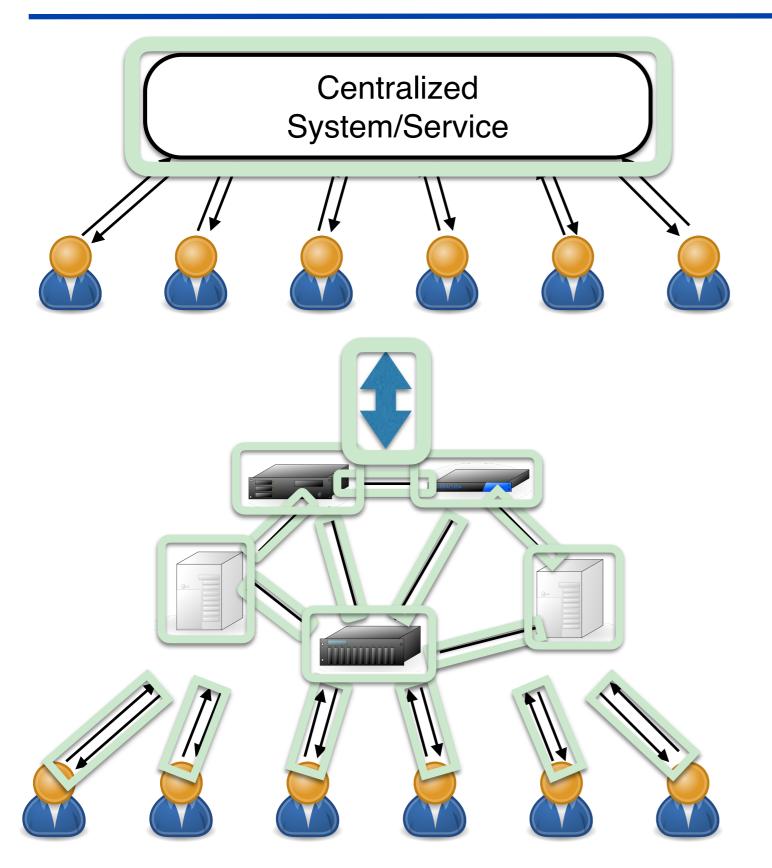
## Blockchain as an Layman



## Roadmap

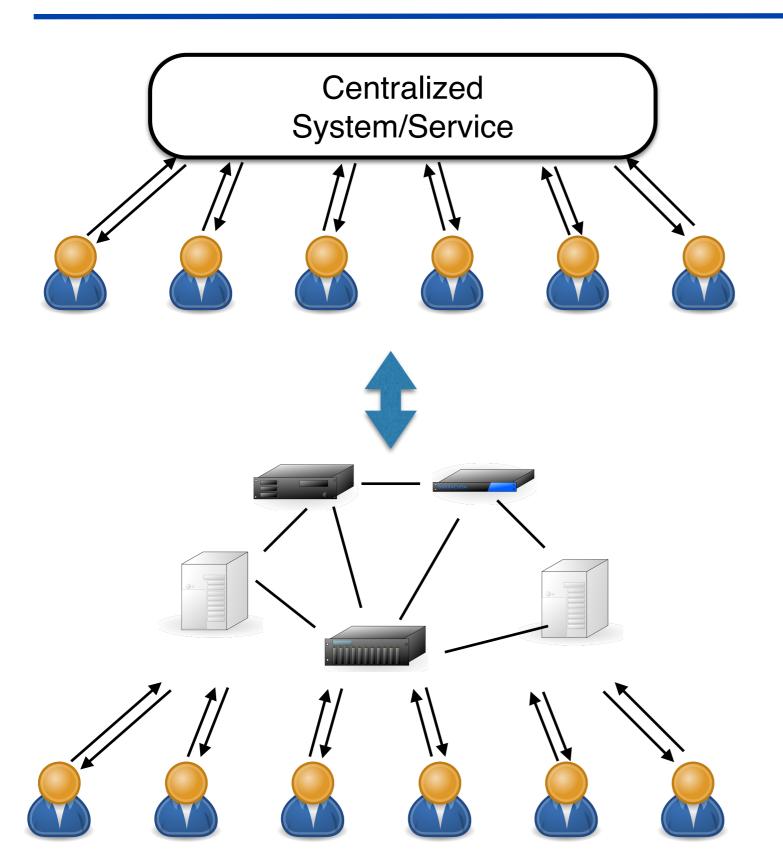
- Part 1: A Principled Design and Analysis of Decentralization with Blockchain
  - Proof of Work, Proof of Stake, and beyond
- Part 2 (in passing...): Economics and blockchain

## **The Decentralization Paradigm**



## Methodology

- 1. Specification
- 2. System Model
- 3. Distributed protocol
- 4. Assumptions and security analysis



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## "Reverse Engineering" Bitcoin

#### Bitcoin: A Peer-to-Peer Electronic Cash System

Satoshi Nakamoto satoshin@gmx.com www.bitcoin.org

Abstract. A purely peer-to-peer version of electronic cash would allow online payments to be sent directly from one party to another without going through a financial institution. Digital signatures provide part of the solution, but the main benefits are lost if a trusted third party is still required to prevent double-spending. We propose a solution to the double-spending problem using a peer-to-peer network. The network timestamps transactions by hashing them into an ongoing chain of hash-based proof-of-work, forming a record that cannot be changed without redoing the proof-of-work. The longest chain not only serves as proof of the sequence of events witnessed, but proof that it came from the largest pool of CPU power. As long as a majority of CPU power is controlled by nodes that are not cooperating to attack the network, they'll generate the longest chain and outpace attackers. The network itself requires minimal structure. Messages are broadcast on a best effort basis, and nodes can leave and rejoin the network at will, accepting the longest proof-of-work chain as proof of what happened while they were gone.

#### 1. Introduction

Commerce on the Internet has come to rely almost exclusively on financial institutions serving as trusted third parties to process electronic payments. While the system works well enough for

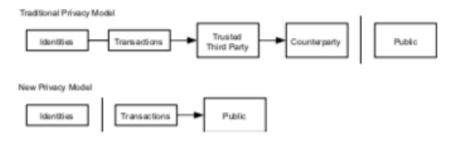
most transactions, it still suffers from the inherent weakness. Completely non-reversible transactions are not really possible, sin avoid mediating disputes. The cost of mediation increases t minimum practical transaction size and cutting off the possibility and there is a broader cost in the loss of ability to make non reversible services. With the possibility of reversal, the need for be wary of their customers, hassling them for more information the A certain percentage of fraud is accepted as unavoidable. These c can be avoided in person by using physical currency, but no mechover a communications channel without a trusted party.

What is needed is an electronic payment system based on cryp allowing any two willing parties to transact directly with each other third party. Transactions that are computationally impractical to from fraud, and routine escrow mechanisms could easily be implicated this paper, we propose a solution to the double-spending problem timestamp server to generate computational proof of the chronolog system is secure as long as honest nodes collectively control to

system is secure as long as honest nodes collectively control more CPU power than any cooperating group of attacker nodes.

#### 10. Privacy

The traditional banking model achieves a level of privacy by limiting access to information to the parties involved and the trusted third party. The necessity to announce all transactions publicly precludes this method, but privacy can still be maintained by breaking the flow of information in another place: by keeping public keys anonymous. The public can see that someone is sending an amount to someone else, but without information linking the transaction to anyone. This is similar to the level of information released by stock exchanges, where the time and size of individual trades, the "tape", is made public, but without telling who the parties were.



As an additional firewall, a new key pair should be used for each transaction to keep them from being linked to a common owner. Some linking is still unavoidable with multi-input transactions, which necessarily reveal that their inputs were owned by the same owner. The risk is that if the owner of a key is revealed, linking could reveal other transactions that belonged to the same owner.

#### 11. Calculations

We consider the scenario of an attacker trying to generate an alternate chain faster than the honest chain. Even if this is accomplished, it does not throw the system open to arbitrary changes, such

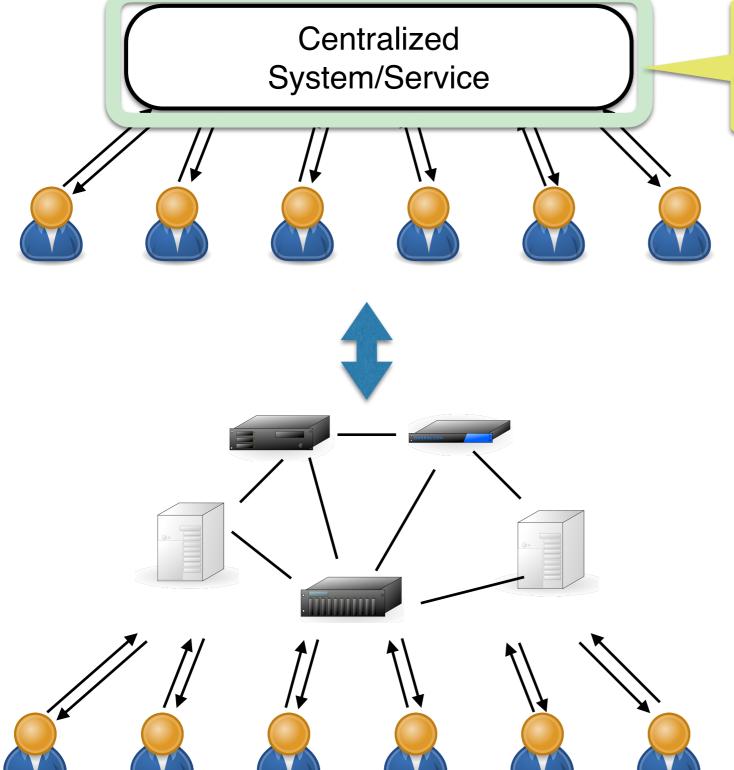
> hin air or taking money that never belonged to the attacker. Nodes are valid transaction as payment, and honest nodes will never accept a block acker can only try to change one of his own transactions to take back

> honest chain and an attacker chain can be characterized as a Binomial ess event is the honest chain being extended by one block, increasing its e event is the attacker's chain being extended by one block, reducing the

> attacker catching up from a given deficit is analogous to a Gambler's a gambler with unlimited credit starts at a deficit and plays potentially an to try to reach breakeven. We can calculate the probability he ever t an attacker ever catches up with the honest chain, as follows [8]:

iest node finds the next block acker finds the next block tacker will ever catch up from z blocks behind

$$q_z = \begin{cases} 1 & \text{if } p \leq q \\ (q/p)^z & \text{if } p > q \end{cases}$$



Blockchain-based Transaction Ledger

## Methodology

- 1. Specification
- 2. System Model
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## **Blockchain-based Transaction Ledger**

## Background: Two common tools from crypto

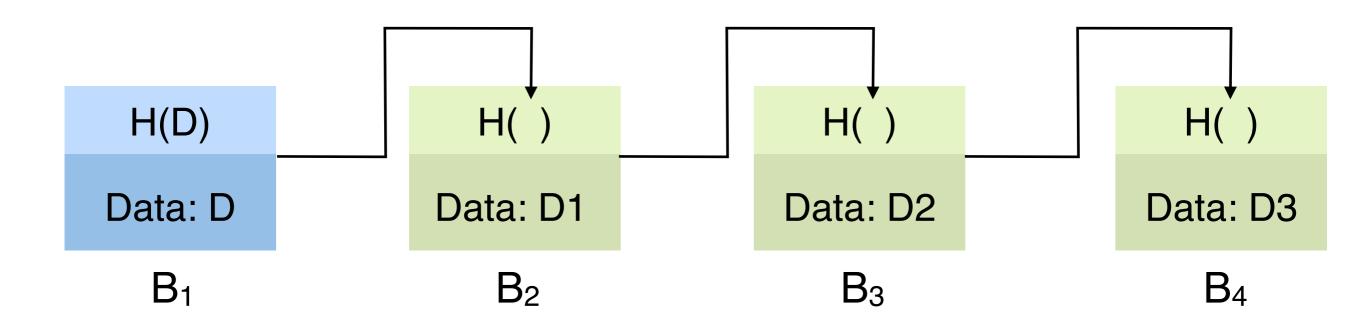
- Hash function H() as a random n-bit function
  - H(x) = 01011...1101 = y
     H(x') = 01000...0101

  - $\bullet$  H(x) = y
- Digital signature sig[ · ]
  - Alice has a secret key sk<sub>Alice</sub>
  - Only Alice can produce signatures sigskalice[ ]
  - Anyone with Alices public key pk<sub>Alice</sub> can verify

**Blockchain-based Transaction Ledger** 

#### **Blockchain-based Transaction Ledger**

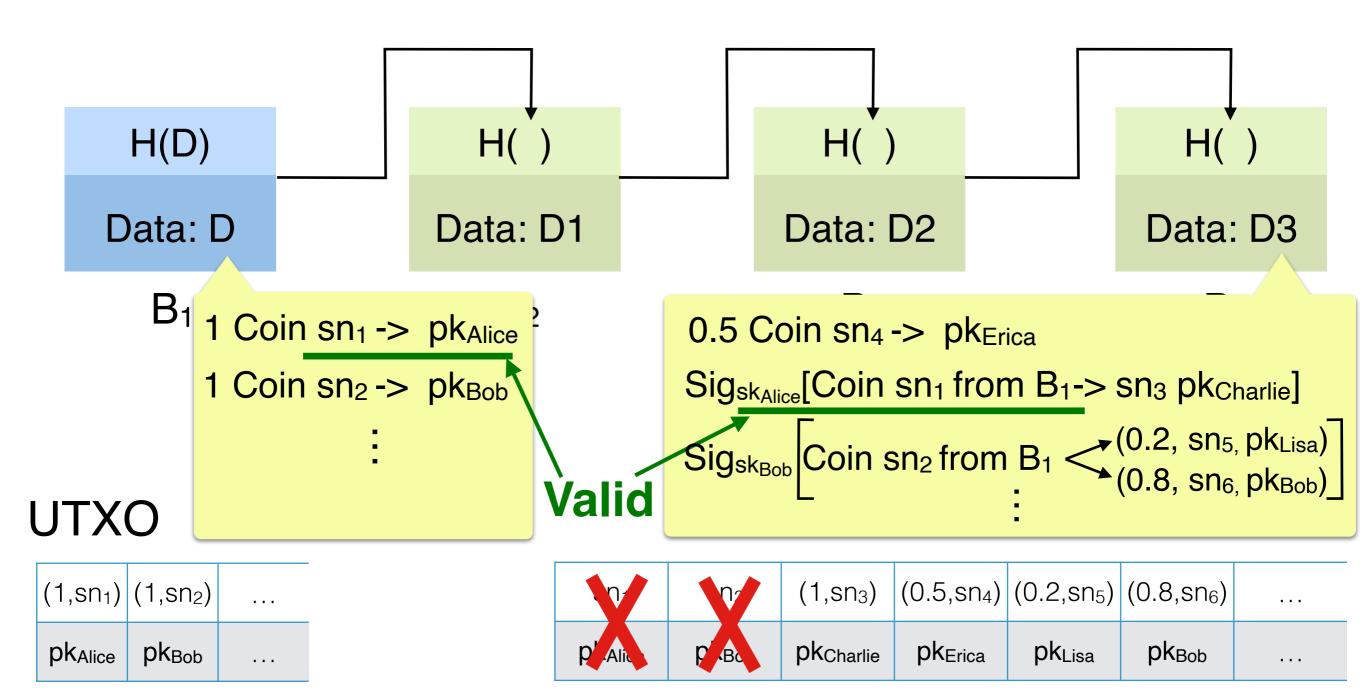
A hash-pointer-based Data Structure



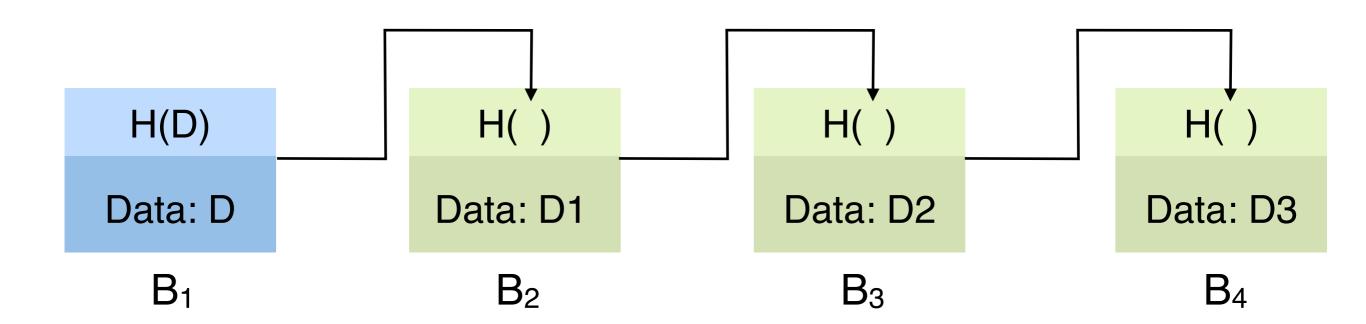
- Causality (aka "time-stamping"): You cannot create block B<sub>i</sub> without knowing block B<sub>i-1</sub>
- Immutability: You cannot modify a block anywhere without changing every follow-up block

## **Blockchain-based Transaction Ledger**

Putting semantics on the blockchain data



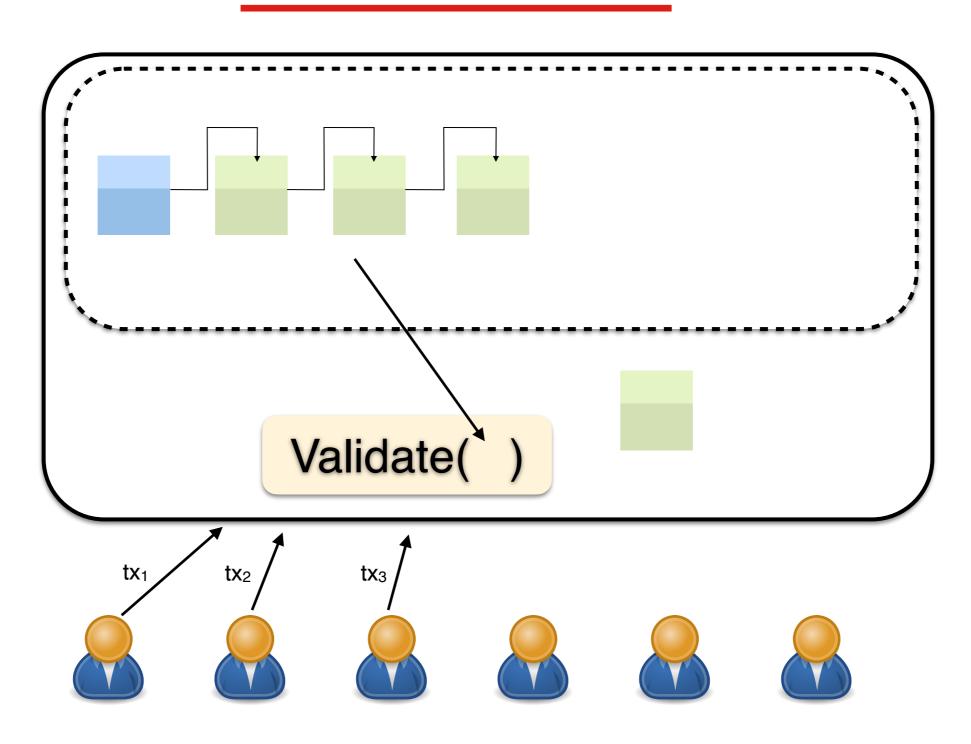
#### **Blockchain-based Transaction Ledger**



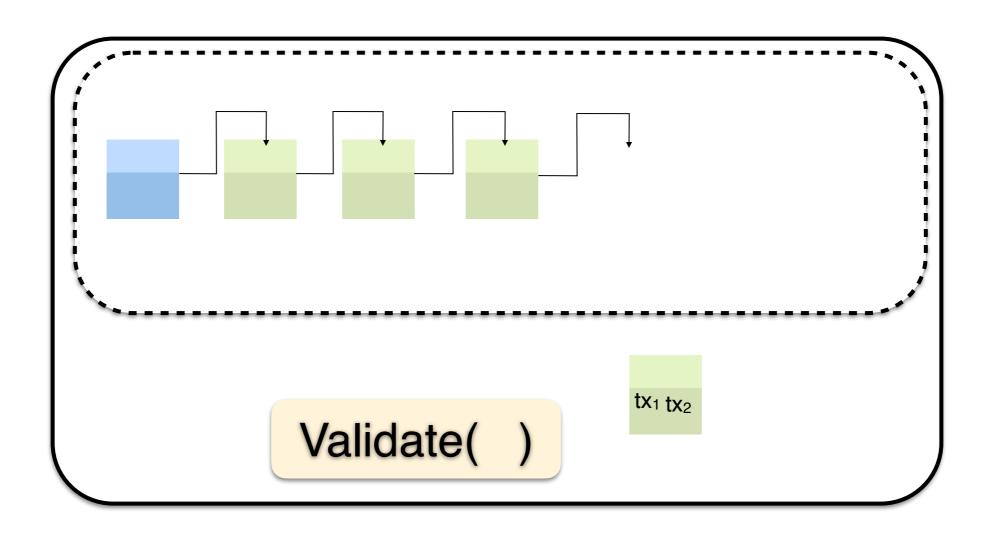
**Blockchain-based Transaction Ledger** 

**Blockchain Ledger** 

## **Blockchain Ledger**



### **Blockchain Ledger**







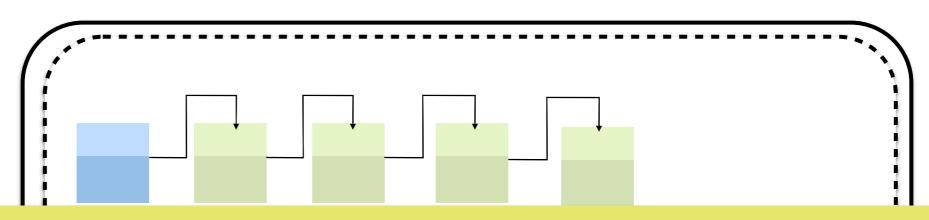






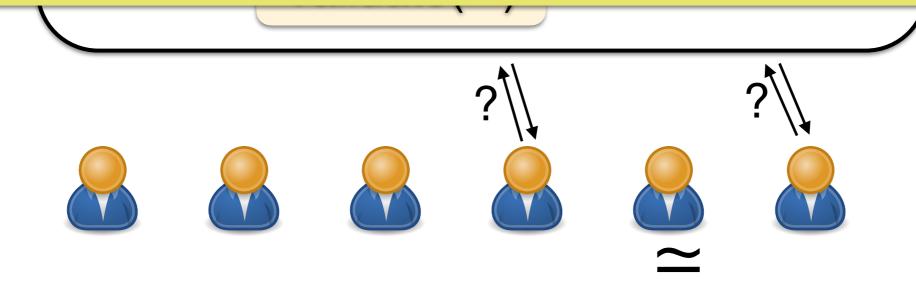


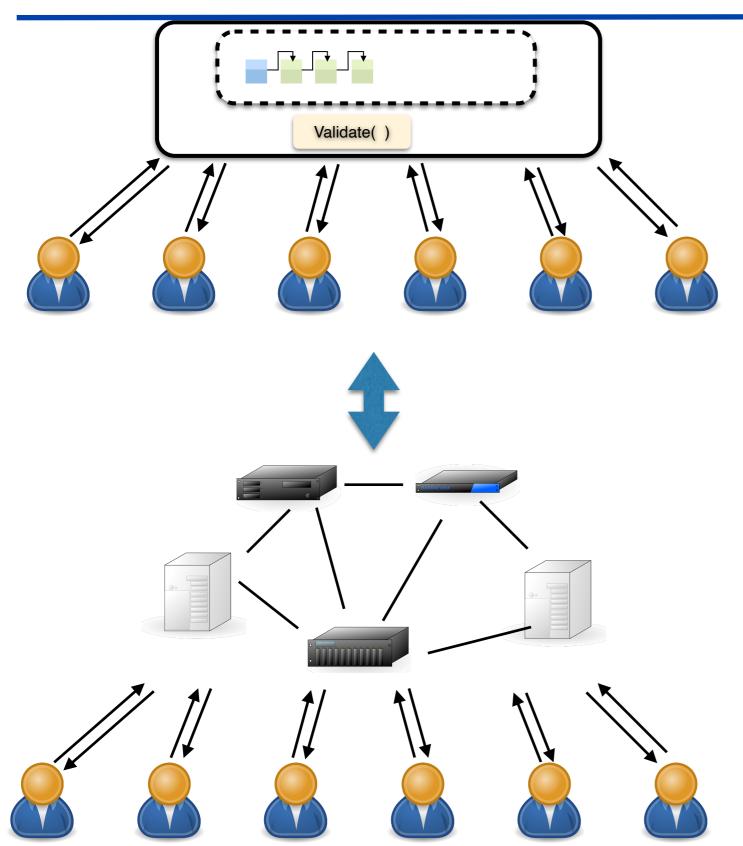
#### **Blockchain Ledger**



[Badertscher-Maurer-Tshudi-Z: CRYPTO17]: Detailed Specification of Bitcoin Ledger:

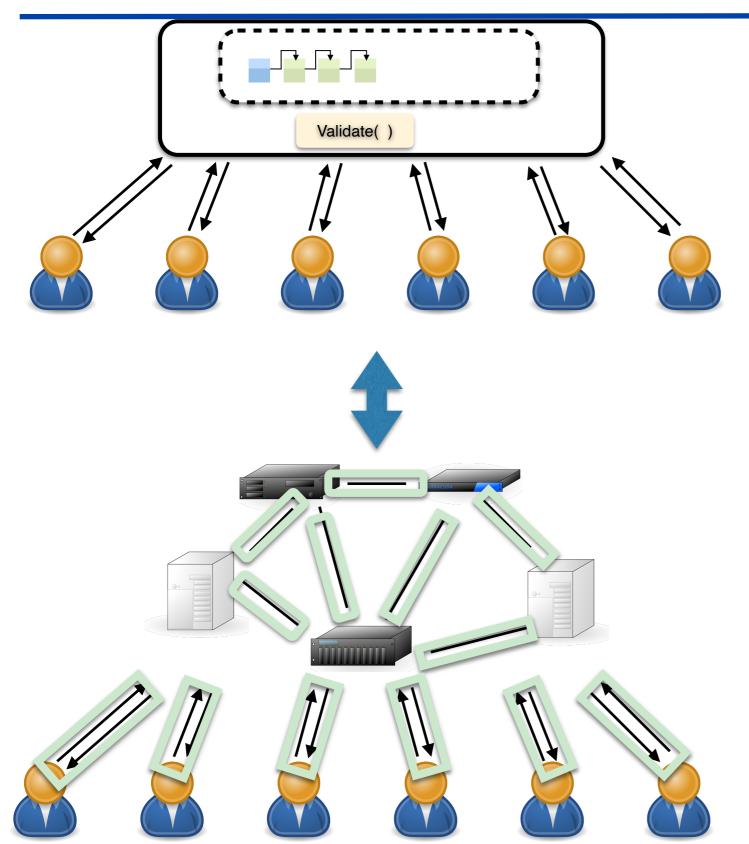
- Outputs might be prefixes of one another
- A transaction might not make it in the immediately next block





## Methodology

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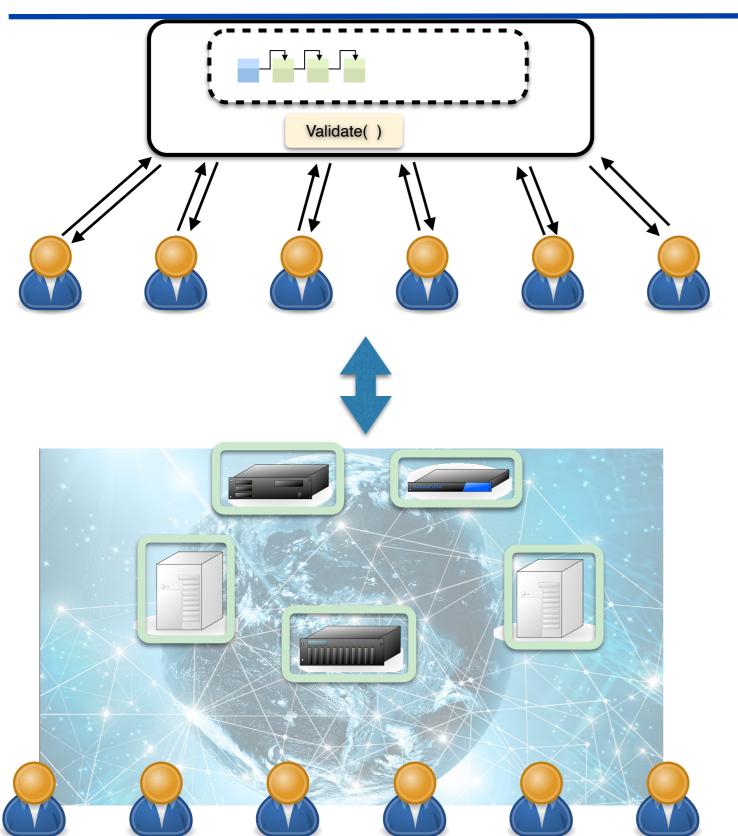
#### **Computing Infrastructure:**

 Simple devices that can hash, do simple computation, store, communicate

#### **Communication:**

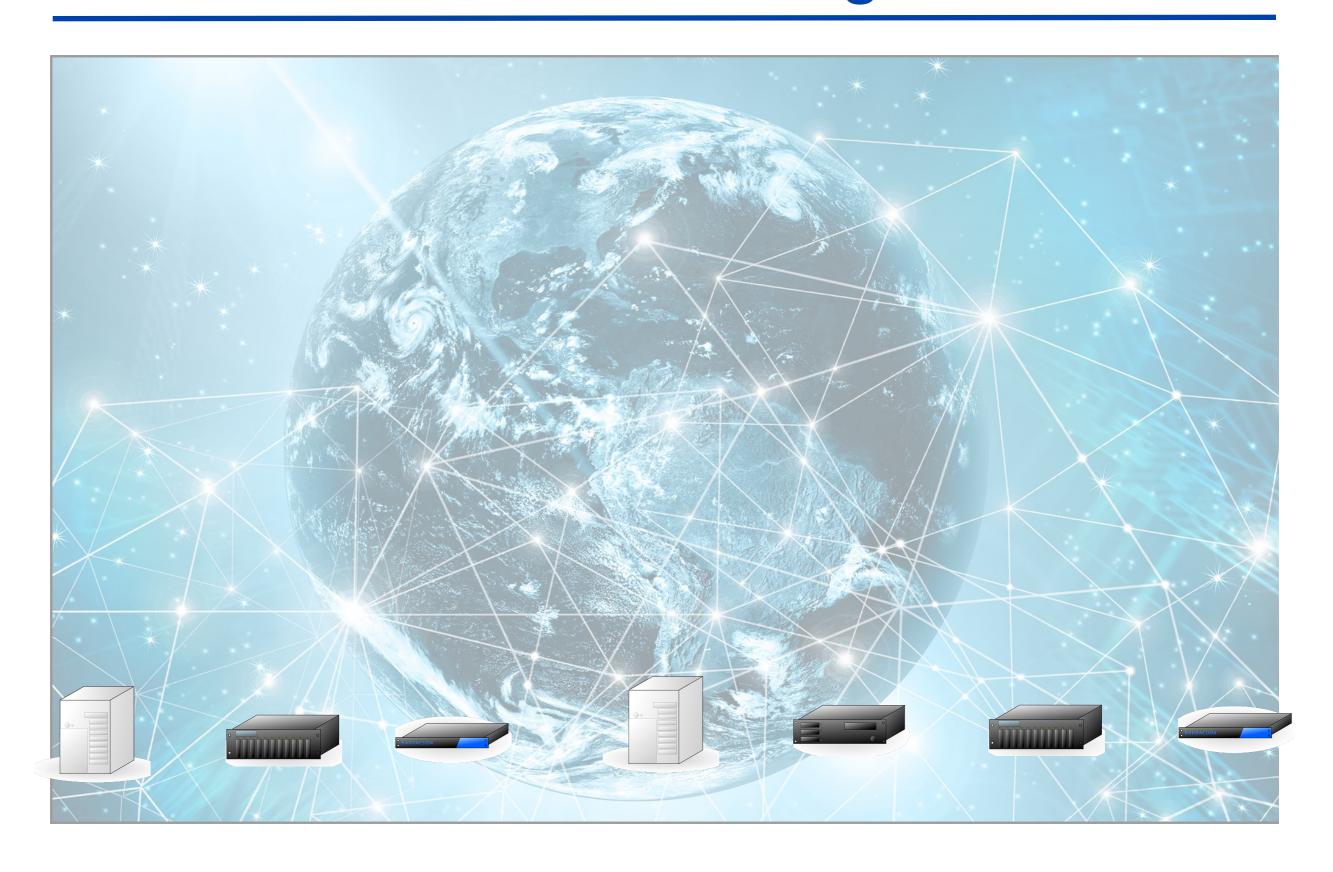
- Challenges:
  - Unknown topology (ad hoc): no direct all to all links
  - Dynamic Participation
- Solution: Gossiping over the Internet



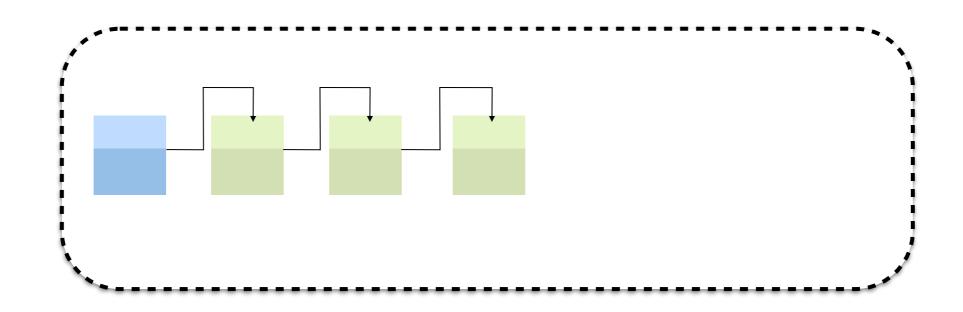


## Methodology

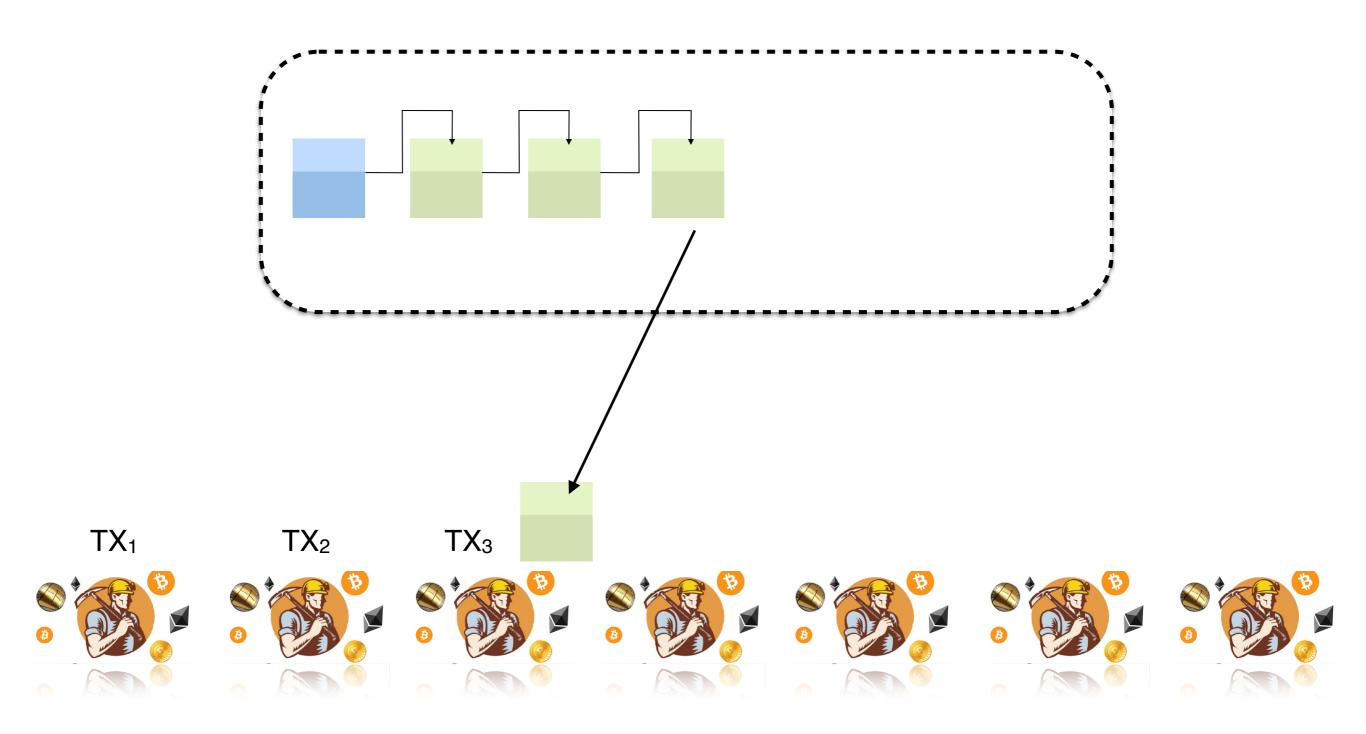
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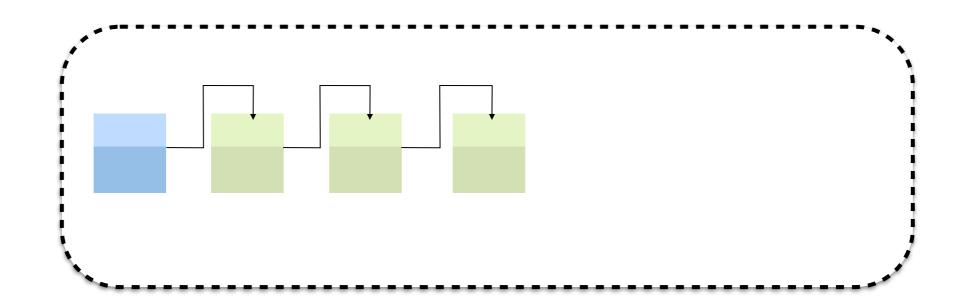




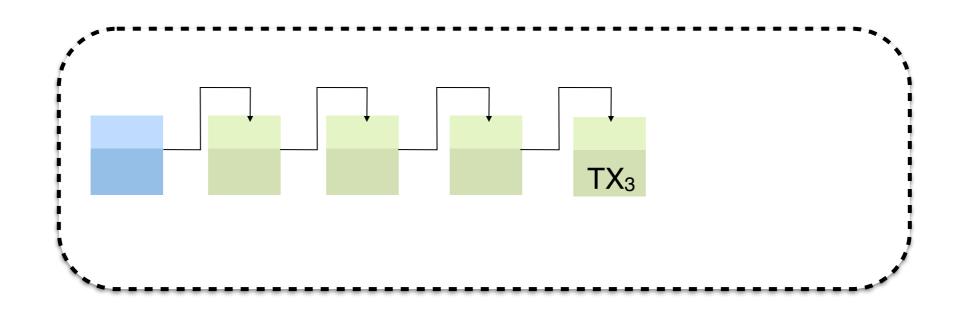














#### Decentralizing a Blockchain Ledger

#### **Key Challenges:**

Who can propose the next block?



- Let's choose randomly
  - From what population?
- How to agree on the blockchain-ledger state?

#### Decentralizing a Blockchain Ledger

#### **Key Challenges:**

Who can propose the next block?





Sybil attack!

SN <sub>1</sub>	sn <sub>2</sub>	(1,sn <sub>3</sub> )	(0.5,sn <sub>4</sub> )	(0.2,sn <sub>5</sub> )	(0.8,sn <sub>6</sub> )	
pk <sub>Alice</sub>	pk <sub>Bob</sub>	pkCharlie	pk <sub>Erica</sub>	pk <sub>Lisa</sub>	pk <sub>Bob</sub>	

How to agree on the blockchain-ledger state?

#### Decentralizing a Blockchain Ledger

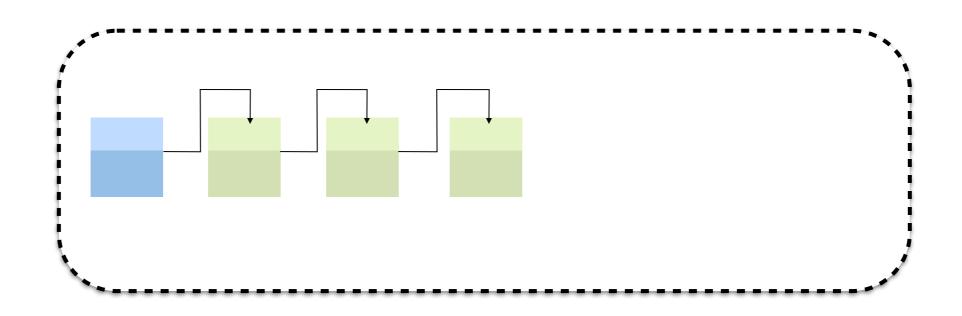
#### **Key Challenges:**

Who can propose the next block?

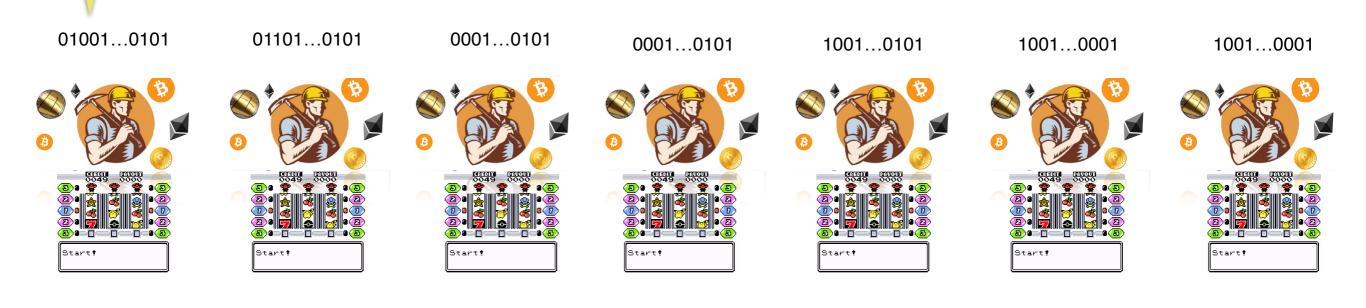


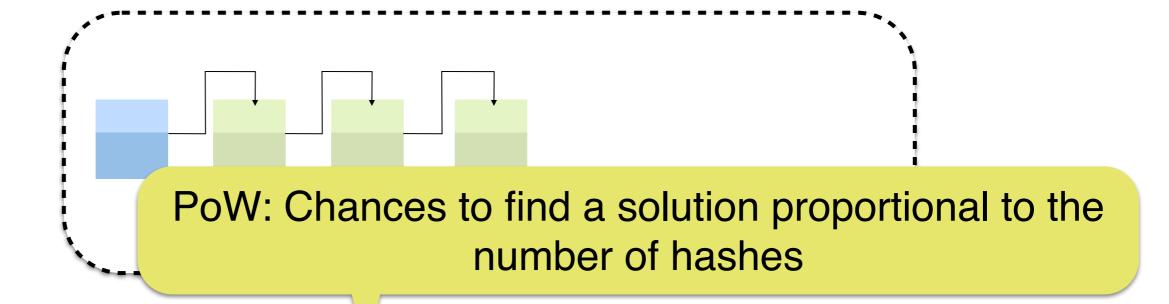
Whoever solves a hash puzzle (proof of work)

How to agree on the blockchain-ledger state?

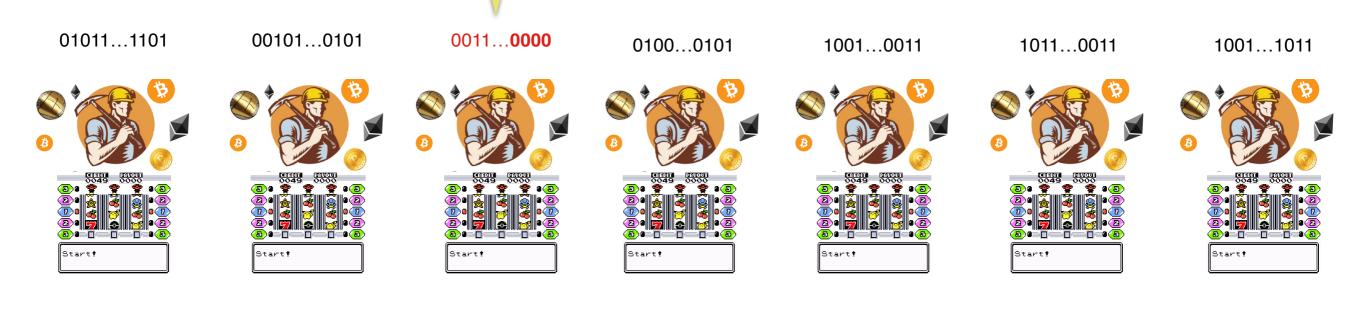


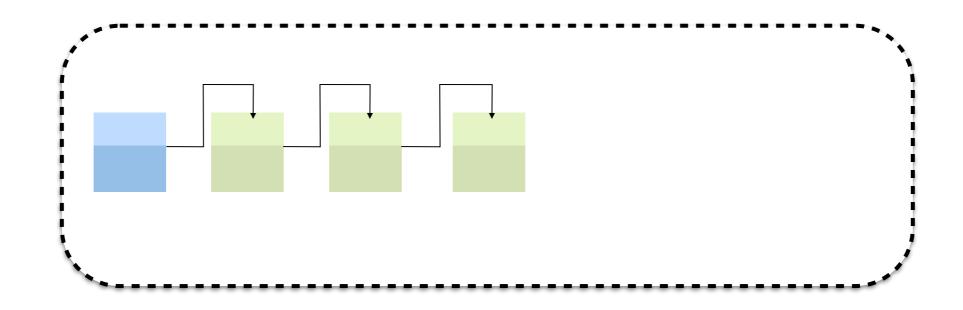
# Hash(Prev. Block, nonce)



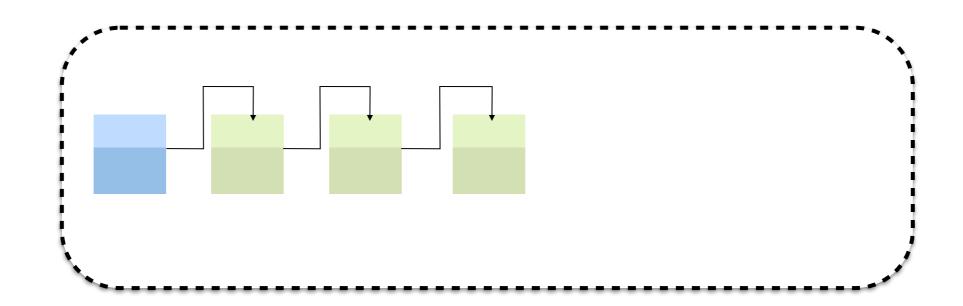


Hash(Prev. Block, nonce)<T















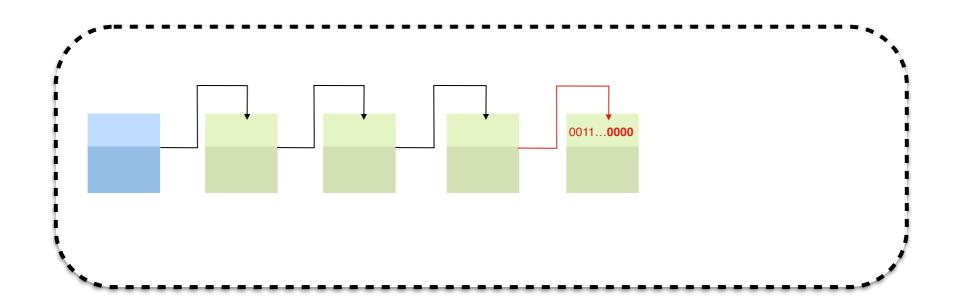
























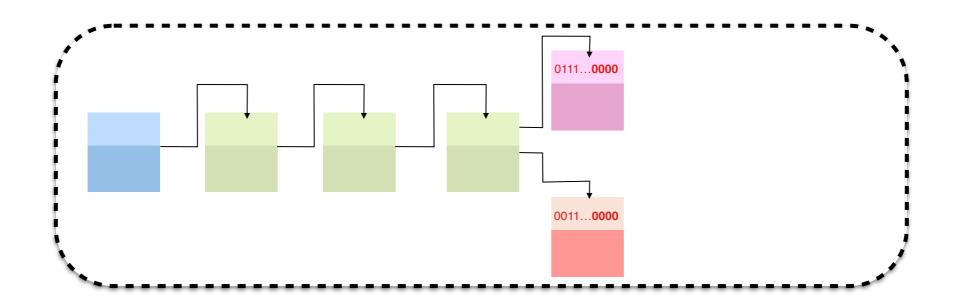


### Decentralizing a Blockchain Ledger

### **Key Challenges:**

- Who can propose the next block?
  - Whoever solves a hash puzzle (proof of work)
  - Chances to win proportional to # of attempts (work)
- How to agree on the blockchain-ledger state?









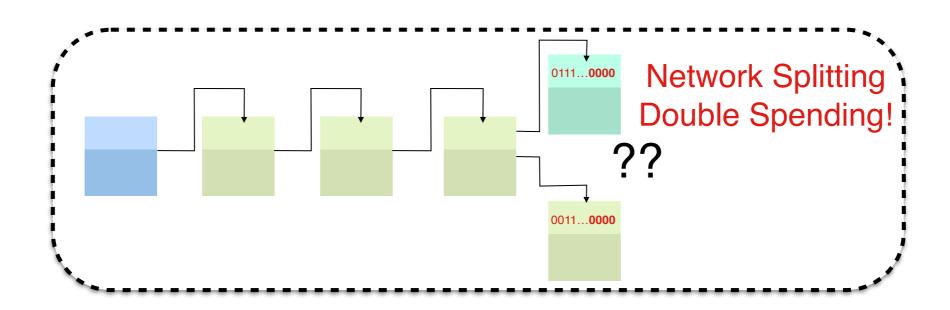














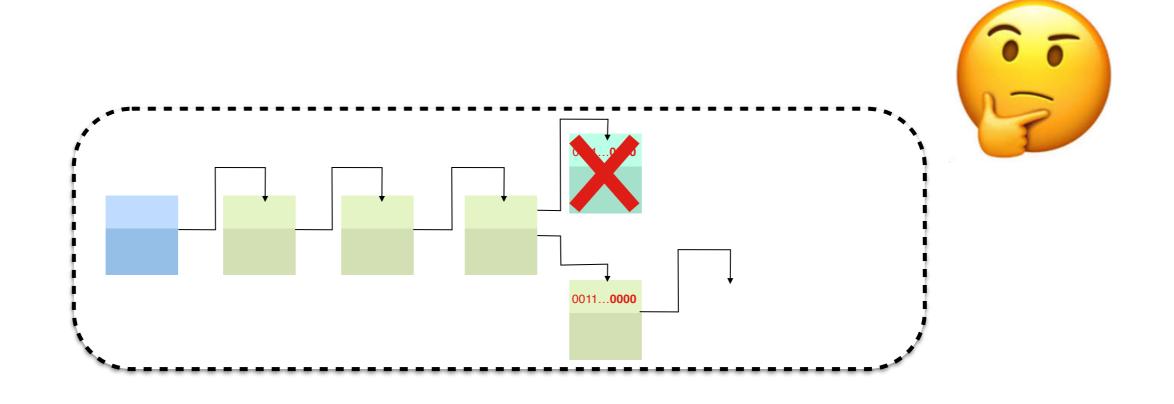
### Decentralizing a Blockchain Ledger

### **Key Challenges:**

- Who can propose the next block?
  - Whoever solves a hash puzzle (proof of work)
  - Chances to win proportional to # of attempts (work)
- How to agree on the blockchain-ledger state?



 Whenever in doubt between two chains adopt the longer (the one with most work)











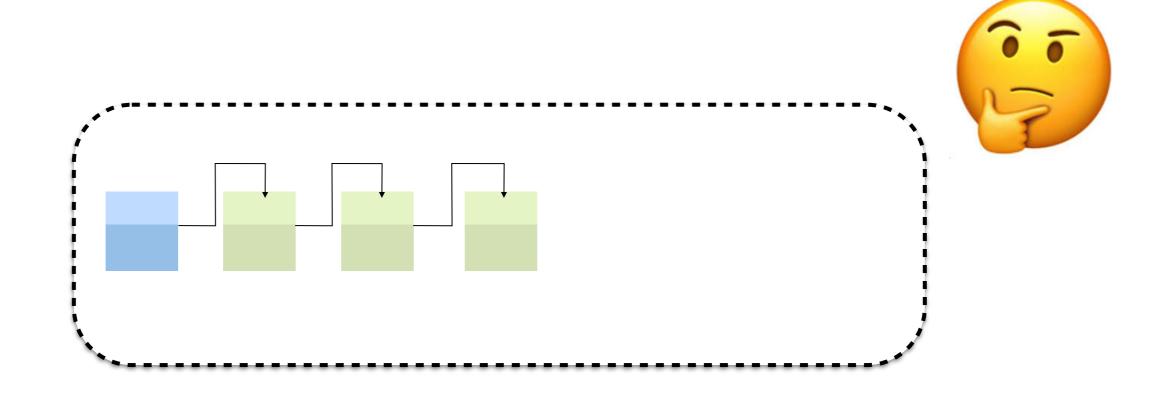












I don't like the last block







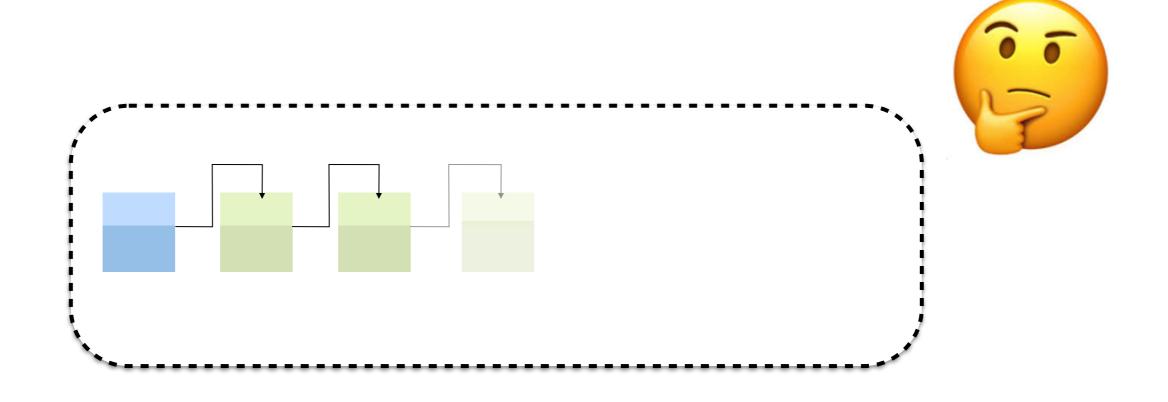
















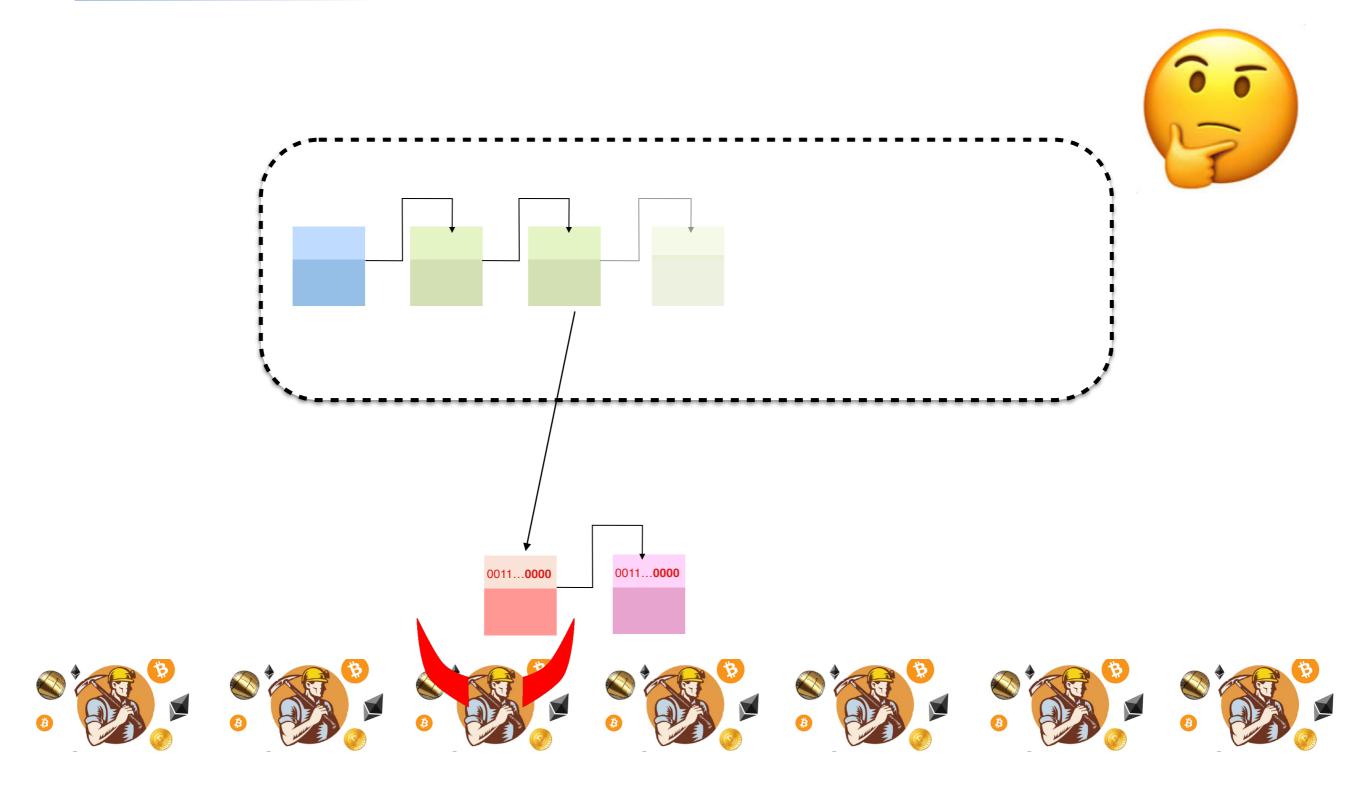


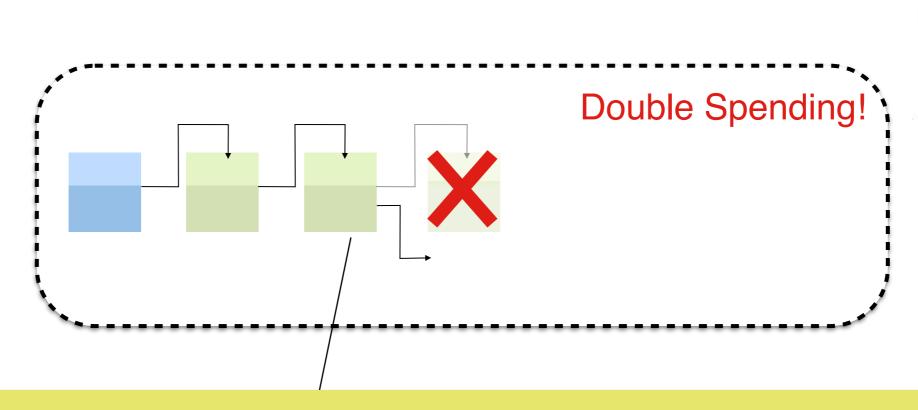












Honest majority of hashing power => Deep enough blocks cannot be orphaned







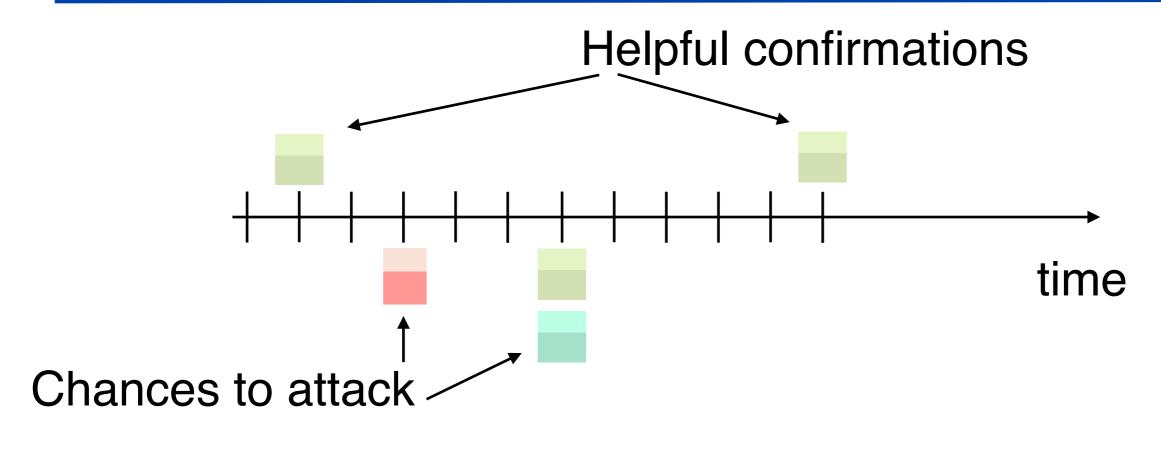


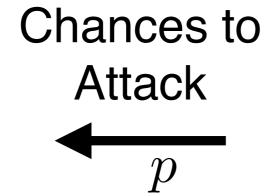




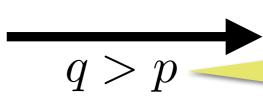








Helpful confirmations



Honest majority of hashing power

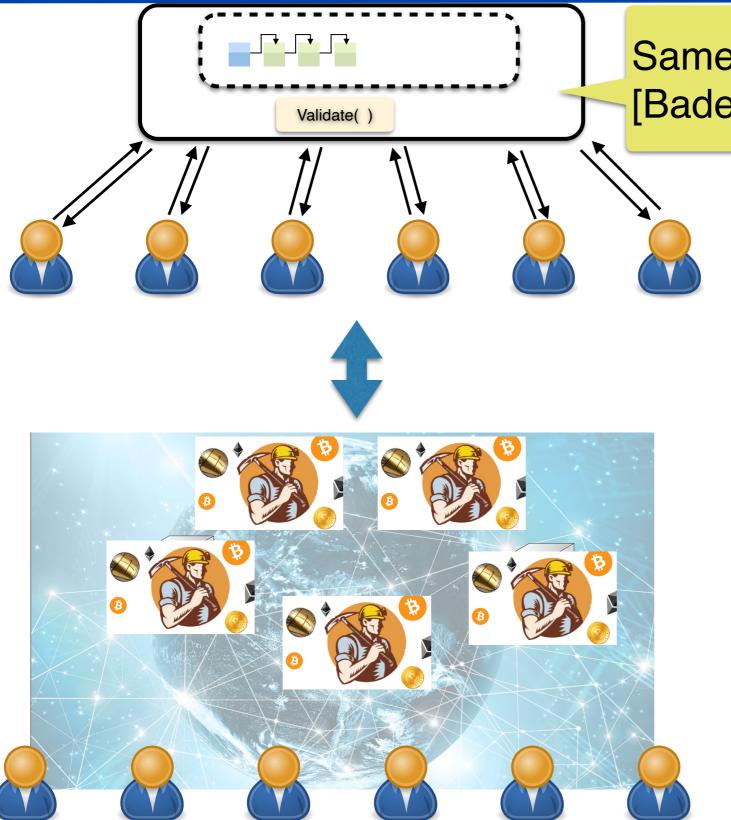
$$\Pr[\operatorname{Left}_k] = Exp(-\Omega(k))$$
block
orphaned

$$\begin{aligned} \Pr[\text{Right}_k] &= 1 - Exp(-\Omega(k)) \\ &\quad \text{block} \\ &\quad \text{confirmed} \end{aligned}$$

### That was over-simplified ...

### Actual security analysis is very delicate:

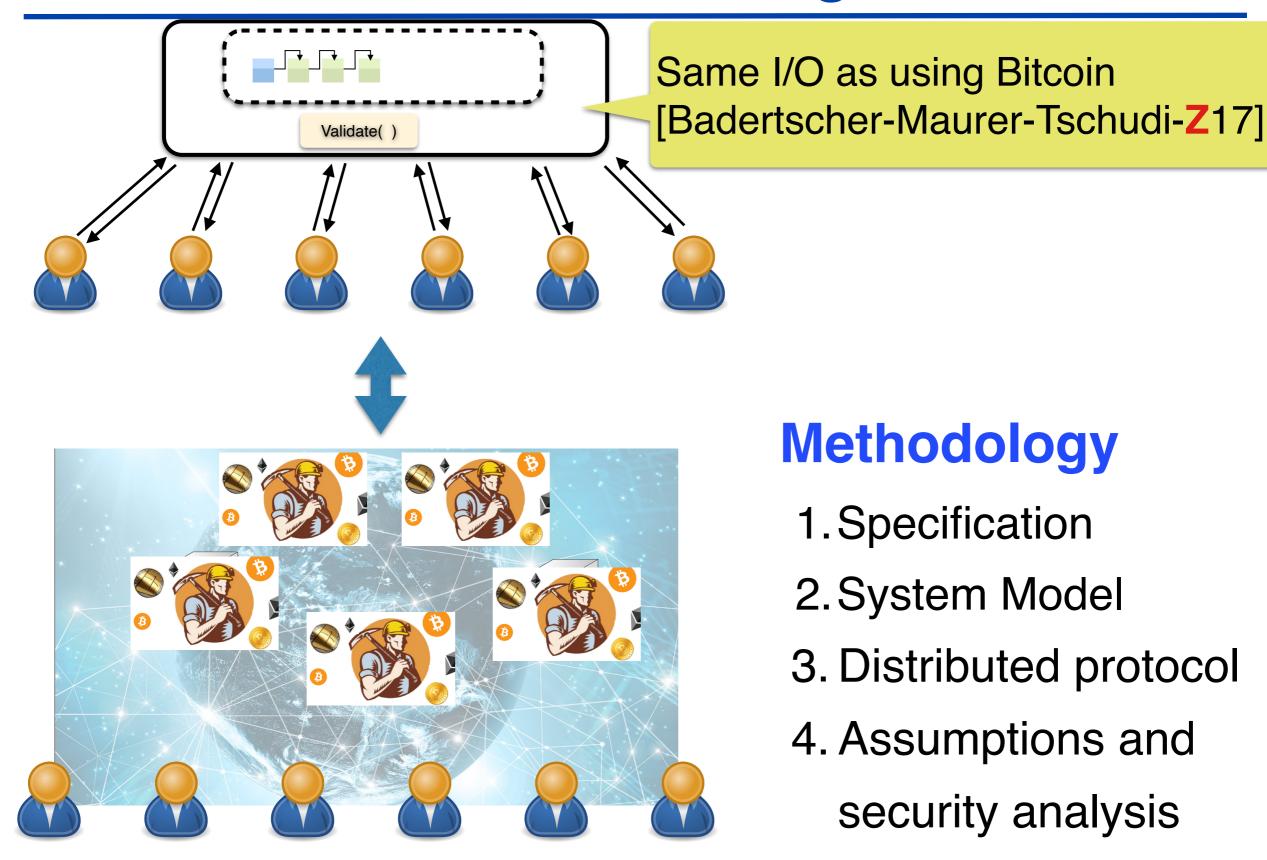
- Block withholding (e.g., selfish mining) reduces honest miners' "effective" hashing
  - When is it optimal to announce?
- Network delay provide opportunities to attacker
  - Buy time to work on a block
  - Temporarily split the network
- Difficulty Readjustment



Same I/O as using Bitcoin [Badertscher-Maurer-Tschudi-Z17]

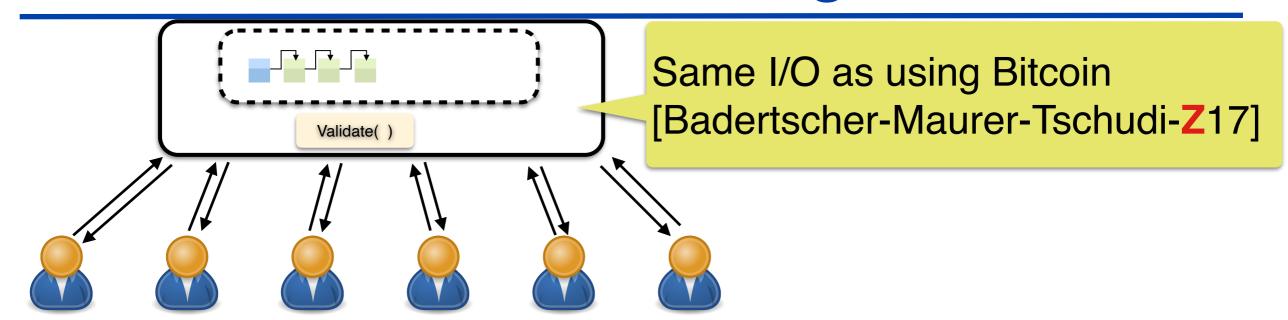
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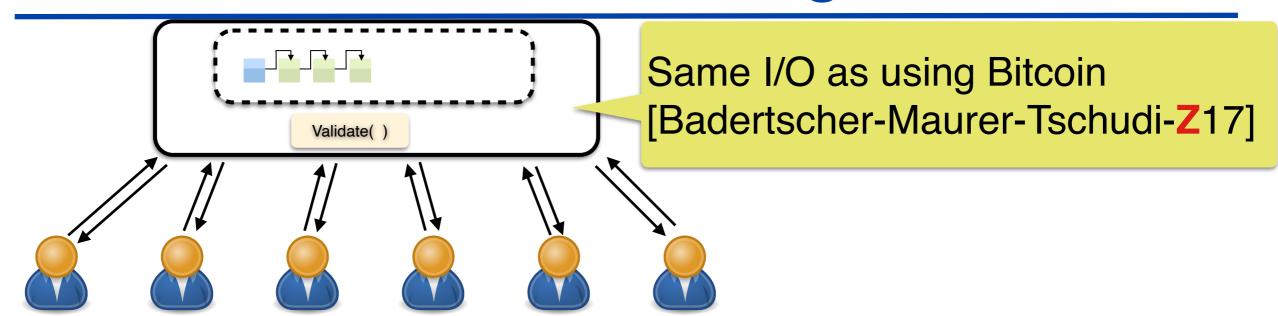
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### Science: Cross-disciplinary research

- Economics
  - Rational analysis of crypto [Badertschere-Lu-Z: EUROCRYPT'20]
  - Incentive-compatible online-poker [Ciampi-Lu-Z: CSF'20]
- AI/ML: Trusted model outsourcing
- Law: Regulation of Cryptocurrencies

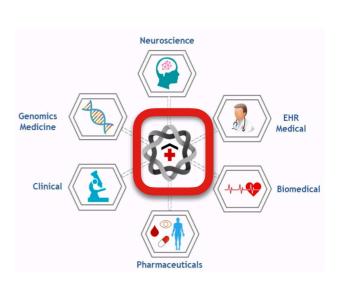


### **Systems Engineering:**

- API to blockchain ("ledger.h")
- Benchmarks specs for different systems







## Proof-of-Work Blockchains (Bitcoin)



### **Economics and Blockchain**

### **Rational Analysis of Blockchains**

The world (Reality)

VS.

Crypto(graphy)



Is Bitcoin (blockchain) secure?

So ... what if the adversary gets majority?



It is secure!!!!! ... **assuming** honest majority of hashing power

not secure ... (or "less secure" if adversarial majority is temporary [BGKRZ20])

#### Economic Robustness (A new type of security statements):

Attacking the assumption and/or the underlying system is irrational

- A good fallback of standard cryptographic security
- Makes modern cryptocurrency DLTs more than "just" a ledger!
- Can improve security/understanding of blockchain-aided protocols

### **Economics and Blockchain**

### **Rational Analysis of Blockchains**

...building on Game-Theoretic Security/ Cryptography [CRYPTO'12, ICALP'12, FOCS'13, PODC15, DISC'15]

- Framework for Economic Robustness of PoW Blockchains
  - [Badertsher-Garay-Maurer-Garay-Tschudi-Z: EUROCRYPT 2018]
    - Rational Crypto Analysis of the bitcoin backbone protocol
- Analysis of 51% Attacks on PoW Blockchains
   [Badertsher-Lu-Z: CRYPTO 2020]
  - Susceptibility and patch

Predicted attack times and cost

Matches forensics result on the ETC attack(s)

$\mathtt{rcost} \; (\mathrm{USD})$	$t  ext{ (days)}$	Cost/day (USD)
\$0.0001	24.0	\$78,084
\$0.0002	10.5	\$156,167
\$0.0003	4.3	\$234,251
\$0.0004	3.2	\$312,334
\$0.0005	2.6	\$390,418
\$0.0006	2.1	\$468,501

DeFi Application: FairMM: Front-running resistant crypto exchanges

[Ciampi-Ishaq-MagdonIsmail-Ostrovsky-Z: CSCML 2022]

A front-running resistant DEX

Feature	FairMM	$\mathbf{Uniswap}$
Front Running Resilience	Yes	No
Gas Price Auctions	No	Yes
Miner Influence	No	Yes
Trade Execuction (seconds)	pprox <b>0.30</b>	$\geq 15$
Average Trade Cost $(K)$	pprox <b>101</b>	$\approx 141^*$
Max Trade Cost $(K)$	pprox <b>101</b>	$pprox 1,316^\dagger$
Max Throughput‡	pprox 475	$\approx 340$

## The Purdue Blockchain Lab (The Pub)



**Blockchain Lab** 

Home Resea

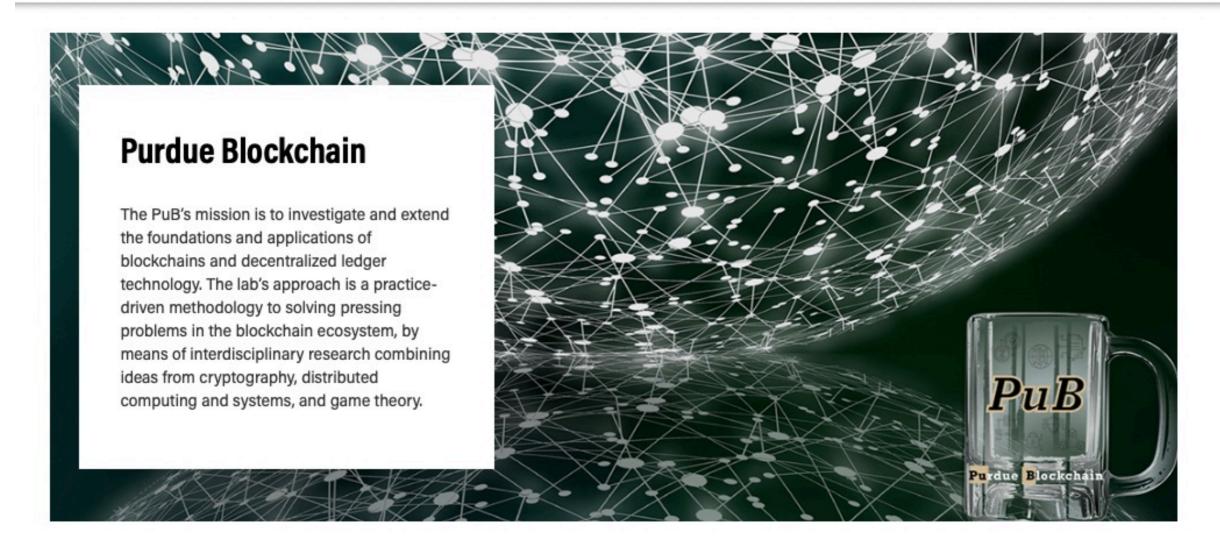
Research

Publications

People Ne

Events

Contact



## The Purdue Blockchain Lab (The Pub)



#### People Lab Director



Vassilis Zikas Associate Professor of Computer Science Vassilis' website

https://www.cs.purdue.edu/blockchain



https://twitter.com/PBlockchainLab

Affiliated Faculty:



Jeremiah Blocki Assistant Professor of Computer Science Jeremiah's website



Christina Garman Assistant Professor of Computer Science Christina's website



Aniket Kate Associate Professor of Computer Science Aniket's website



Alex Psomas Assistant Professor of Computer Science Alex's website

Thank you!

### **MEGA-ACE: Purdue Partners with NTUA**

